

# Combinatorial Optimization of Matrix-Vector Multiplication

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**Abstract:** Work by Kirby, et al. (2006) showed that combinatorial optimization of matrix/vector multiplication could lead to faster evaluation of finite element stiffness matrices. Using relationships between rows, an efficient set of operations can be generated to perform matrix-vector multiplication. My improved graph model of this problem solves this combinatorial optimization problem optimally for binary row relationships. I extend the representation by using hypergraphs to model more complicated row relationships, expressing an n-row relationship with an n-vertex hyperedge. My initial greedy algorithm for this hypergraph model has yielded significantly better results than the graph model for many matrices.

## Motivation

- Reducing redundant operations in building stiffness matrices
  - Generate code to optimize construction of local stiffness matrices
  - Many local stiffness matrices built
  - Reuse optimized code when problem is rerun
- Finite element "Compilers"
  - FIAT (automates generation of FEs)
  - FFC (variational forms -> code for evaluation)
- Following on work by Kirby, et al., Texas Tech, University of Chicago
  - Optimization of FFC generated code
  - FErari

## Matrix-Vector Multiplication

Kirby, et al. showed that each element of FE local stiffness matrix can be calculated by the following Frobenius product:

$$S_{i,j}^e = K_{i,j} : G^e = \sum_m \sum_n G_{m,n}^e K_{i,j,m,n}$$

Flattening the stiffness matrix into a vector, we get a matrix-vector multiplication operation ( $y=Ax$ ) with the following inner products for each row:

$$y_{ni+j} = S_{i,j}^e = A_{(ni+j,*)} x$$

For the 2D Laplace equation, for example, we get the inner products of the following vectors:

$$A_{(ni+j,*)}^T = \begin{pmatrix} \frac{\partial \phi_i}{\partial r} & \frac{\partial \phi_j}{\partial r} \\ \frac{\partial \phi_i}{\partial s} & \frac{\partial \phi_j}{\partial s} \\ \frac{\partial \phi_i}{\partial x} & \frac{\partial \phi_j}{\partial x} \\ \frac{\partial \phi_i}{\partial y} & \frac{\partial \phi_j}{\partial y} \end{pmatrix} \begin{pmatrix} \epsilon \\ \epsilon \\ \epsilon \\ \epsilon \end{pmatrix} \quad x = \det(J) \begin{pmatrix} \frac{\partial x}{\partial r} \frac{\partial x}{\partial s} + \frac{\partial x}{\partial r} \frac{\partial x}{\partial y} \\ \frac{\partial x}{\partial x} \frac{\partial x}{\partial x} + \frac{\partial x}{\partial y} \frac{\partial x}{\partial y} \\ \frac{\partial x}{\partial x} \frac{\partial x}{\partial x} + \frac{\partial x}{\partial y} \frac{\partial x}{\partial y} \\ \frac{\partial x}{\partial x} \frac{\partial x}{\partial x} + \frac{\partial x}{\partial y} \frac{\partial x}{\partial y} \end{pmatrix}$$

↑ Element independent      ↑ Element dependent

## Optimization Problem

**Objective:** To minimize the number of multiply add pairs (MAPs) in matrix/vector multiplication

e.g.  $r_2 = \alpha r_1 \Rightarrow y_2 = \alpha y_1$

$$\begin{bmatrix} y_1 \\ y_2 \\ \vdots \\ y_n \end{bmatrix} = \begin{bmatrix} r_1^T \\ \alpha r_1^T \\ \vdots \\ r_m^T \end{bmatrix} x = \begin{bmatrix} r_1^T x \\ \alpha y_1 \\ \vdots \\ r_m^T x \end{bmatrix}$$

## Possible Optimizations

**Number of Nonzeros**

$$\begin{aligned} y_3 &= 0 && \boxed{0 \text{ MAPs}} \\ y_1 &= 2x_1 && \boxed{1 \text{ MAP}} \\ y_2 &= 2x_1 + 2x_2 && \boxed{2 \text{ MAPs}} \\ y_4 &= 2x_1 + 2x_2 + 2x_3 && \boxed{3 \text{ MAPs}} \end{aligned}$$

**Scalar Multiples**

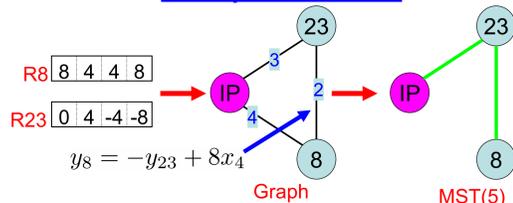
$$\begin{aligned} y_4 &= y_2 && \boxed{0 \text{ MAPs}} \\ y_2 &= 4y_1 && \boxed{1 \text{ MAP}} \\ y_2 &= 4y_4 + 4x_4 && \boxed{2 \text{ MAPs}} \end{aligned}$$

**Linear Combinations**

$$y_2 = 4y_3 + 4y_4 \quad \boxed{2 \text{ MAPs}}$$

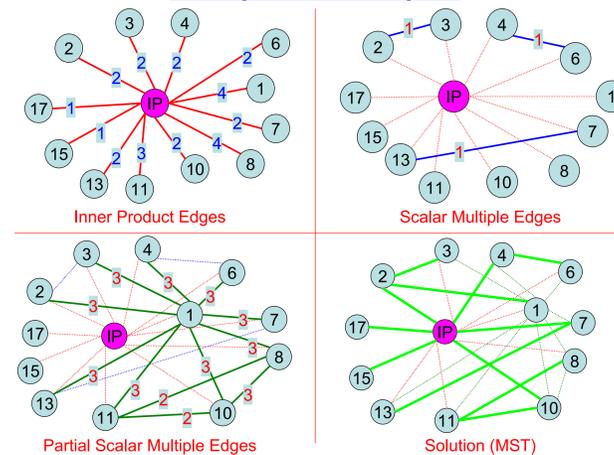
Partial Scalar Multiple →

## Graph Model



- Relationships between rows represented by edges between corresponding vertices
- Additional edges connect each vertex to IP vertex, representing inner product
  - FErari does not include this relationship in graph model
- Edge weights are MAP cost for particular operation
- Edges in minimum spanning tree (MST) represent operations to optimally compute (with these operations) matrix-vector multiplication
- In my work, I used Prim's algorithm to calculate MST

## Graph Example



## Graph Model Results

### 2D Laplace FErari Matrices, Matvec Costs

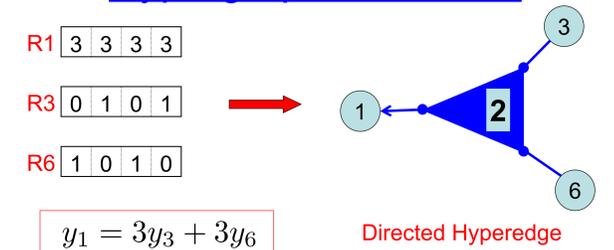
Order	Unoptimized MAPs	Nonzero MAPs	FErari MAPs	My MAPs
1	18	10	7	7
2	63	34	15	14*
3	165	108	45	43
4	360	292	176	152
5	693	589	443	366
6	1218	1070	867	678

### 3D Laplace FErari Matrices, Matvec Costs

Order	Unoptimized MAPs	Nonzero MAPs	FErari MAPs	My MAPs
1	60	21	-	17
2	330	177	94	79
3	1260	789	362	342
4	3780	2586	1118	1049
5	9576	7125	-	3567

\*optimal

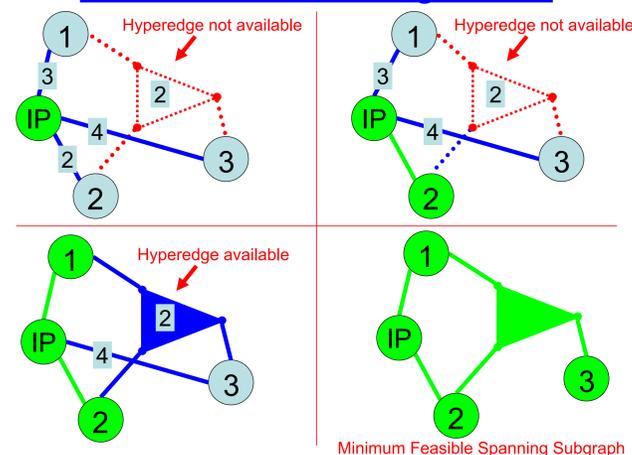
## Hypergraph Extension



## First Pass: Greedy Algorithm

- Modified Prim's algorithm to include hyperedges
- Polynomial time algorithm
- Solution no longer a tree
- No guarantee of optimum solution for hypergraph

## Modified Prim's Algorithm



## Hypergraph Model Results

### 2D Laplace FErari Matrices, Matvec Costs

Order	Unoptimized MAPs	Nonzero MAPs	Graph MAPs	HGraph MAPs
1	18	10	7	6*
2	63	34	14*	14*
3	165	108	43	43
4	360	292	152	150
5	693	589	366	363
6	1218	1070	678	678

### 3D Laplace FErari Matrices, Matvec Costs

Order	Unoptimized MAPs	Nonzero MAPs	FErari Geom. MAPs	My Graph MAPs	My HGraph MAPs
1	60	21	-	17	14*
2	330	177	105	79	68
3	1260	789	327	342	296
4	3780	2586	1045	1049	833
5	9576	7125	-	3567	3250

\*optimal

## Matrix-Vector Multiplication Timings

### 2D Laplace FErari Matrices

Order	Original Time (s)	Optimized Time (s)
1	$2.16 \times 10^{-3}$	$1.80 \times 10^{-4}$
2	$7.74 \times 10^{-3}$	$6.95 \times 10^{-4}$
3	$1.94 \times 10^{-2}$	$1.87 \times 10^{-3}$
4	$4.23 \times 10^{-2}$	$5.20 \times 10^{-3}$
5	$8.09 \times 10^{-2}$	$1.02 \times 10^{-2}$
6	$1.53 \times 10^{-1}$	$1.85 \times 10^{-2}$

### 3D Laplace FErari Matrices

Order	Original Time (s)	Optimized Time (s)
1	$7.09 \times 10^{-3}$	$3.65 \times 10^{-4}$
2	$3.70 \times 10^{-2}$	$1.68 \times 10^{-3}$
3	$1.72 \times 10^{-1}$	$9.62 \times 10^{-3}$
4	$5.96 \times 10^{-1}$	$3.15 \times 10^{-2}$

- 10,000 matrix vector multiplications per timing

## Future Work

- Larger size hyperedges
  - Have implemented 3 and 4 vertex hyperedges
  - Higher degree perhaps useful for 3D FE problems
- More nearly optimal and optimal solution methods

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## References

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